

Wi-Fi Technology Fundamentals

Module-3 **WLAN MAC Layer** Session-3d Data Transfer and Aggregation





Last Session Recap.....



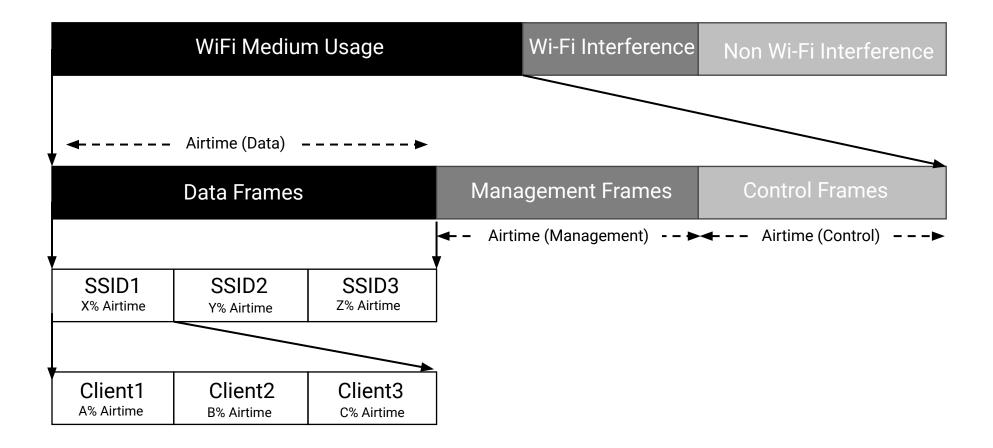
Module-3 WLAN MAC Layer Session-3c Carrier Sense and Medium Access

- ✓ CSMA/CD vs CSMA/CA
- PCF and DCF
- Network Allocation Vector
- Random Backoff
- ✔ Interframe Space
- Contention Window
- EDCA Parameter Set
- ✔ 802.11e QoS
- Basic Demo

Wi-Fi Airtime

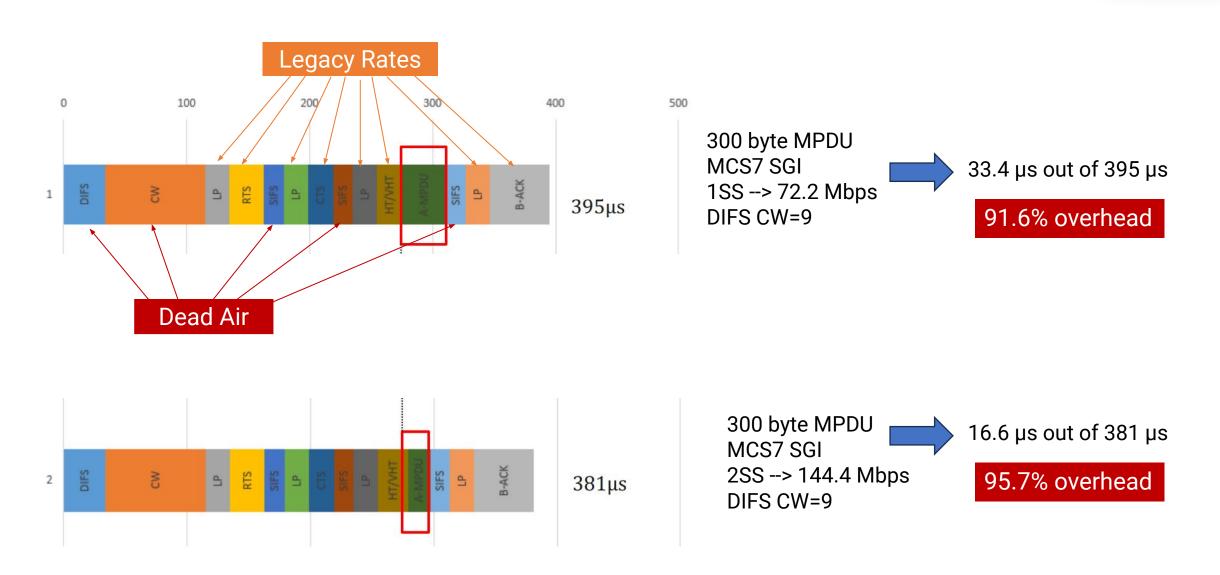


- Airtime taken by actual Wi-Fi Tx/Rx within a cell and then Wi-Fi interference from neighboring APs and also non Wi-Fi interference.
- Of the actual Tx/Rx in the cell, a small percentage is used for Data Frames, rest goes to Management and Control Frames.
- Of the Data frame, airtime is then split between various WiFi clients.



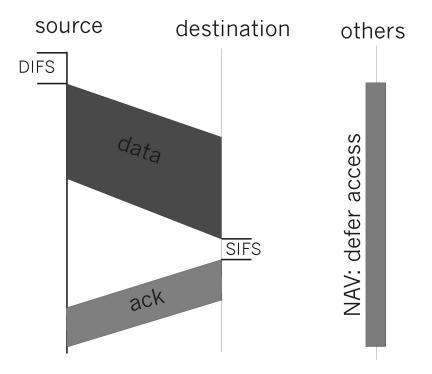
Wi-Fi Protocol and Medium Access Overhead





Data Frame Transmission Overhead

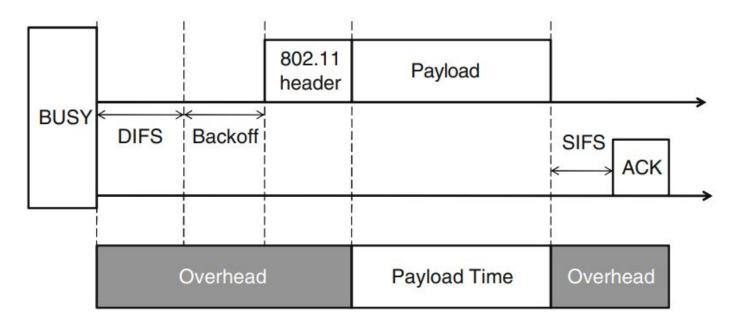




Overhead comes from:

- Interframe Spaces
- Random Backoff
- PHY and MAC Headers
- Smaller Frame Sizes
- ACKs
- Retries

| | Frame Size (Bytes) | | | | | | |
|----------------|--------------------|----------------------|------------|----------------------|--|--|--|
| IEEE 802.11a/g | 10 | 0 | 1500 | | | | |
| | Throughput | % Loss efficiency | Throughput | % Loss efficiency | | | |
| 6 Mbps | 2.164 | 63.93 | 5.372 | 10.46 | | | |
| 9 Mbps | 2.687 | 70.14 | 7.784 | 13.51 | | | |
| 12 Mbps | 3.011 | 74.91 | 10.019 | 16.51 | | | |
| 18 Mbps | 3.483 | 80.65 | 14.123 | 21.54 | | | |
| 24 Mbps | 3.744 | 84.40 | 17.603 | 26.65 | | | |
| 36 Mbps | 4.047 | 88.76 | 23.543 | 34.60 | | | |
| 48 Mbps | 4.217 | 91.21 | 28.189 | 41.27 | | | |
| 54 Mbps | 4.217 | 92.19 | 30.480 | 43.56 | | | |



Data Re-Transmission

DIFS

Transmitter

Receiver

Random

Backoff

Data

SIFS

ACK Timeout

- Retries can exponentially increase the amount of overhead
- Every retry will increase backoff and will reduce PHY rates that will be increase the airtime consumed for transmission of the same frame.

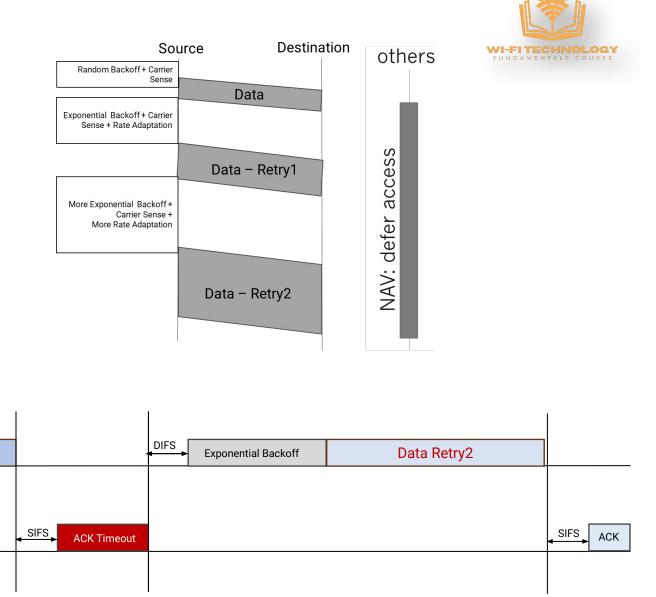
DIFS

Exponential

Backoff

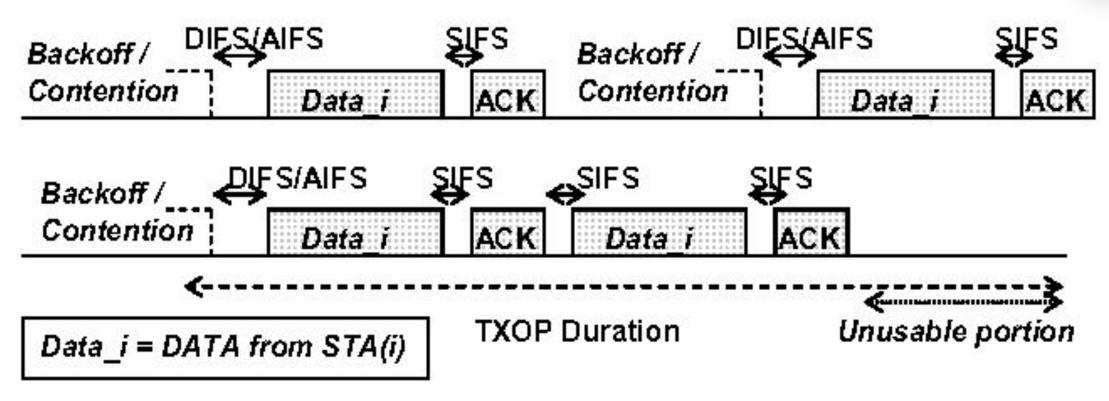
Data Retry1

Extra Overhead



Reducing overhead using TXOP





TXOP will allow device to transmit multiple frames without doing backoff/contention for each frame. This allows for a substantial decrease in overhead.

The Airtime used by Frames

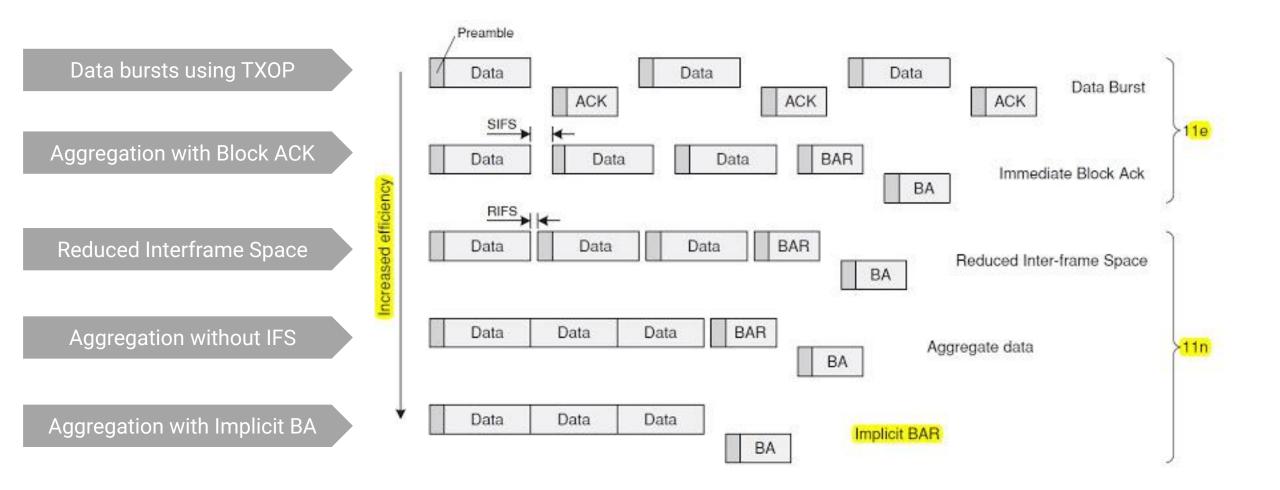


| Modulation | Coding | Nss1 SGI 160Mhz | Airtime for 1500 Byte Frame | Number of Frames in same airtime | 185 μs Timeslot |
|------------|--------|--------------------|-----------------------------------|---|-----------------|
| BPSK | 1/2 | 65.0 | 184.62 | 1 | |
| QPSK | 1/2 | 130.0 | 92.31 | 2 | |
| QPSK | 3/4 | 195.0 | 61.54 | 3 | |
| 16-QAM | 1/2 | 260.0 | 46.15 | 4 | |
| 16-QAM | 3/4 | 390.0 | 30.77 | 6 | |
| 64-QAM | 2/3 | 520.0 | 23.08 | 8 | |
| 64-QAM | 3/4 | 585.0 | 20.51 | 9 | |
| 64-QAM | 5/6 | 650.0 | 18.46 | 10 | |

The higher the data rates the more frames that can be transmitted in the same amount of time.

Methods used to increase thought put and reduce overhead





Concept of Aggregation



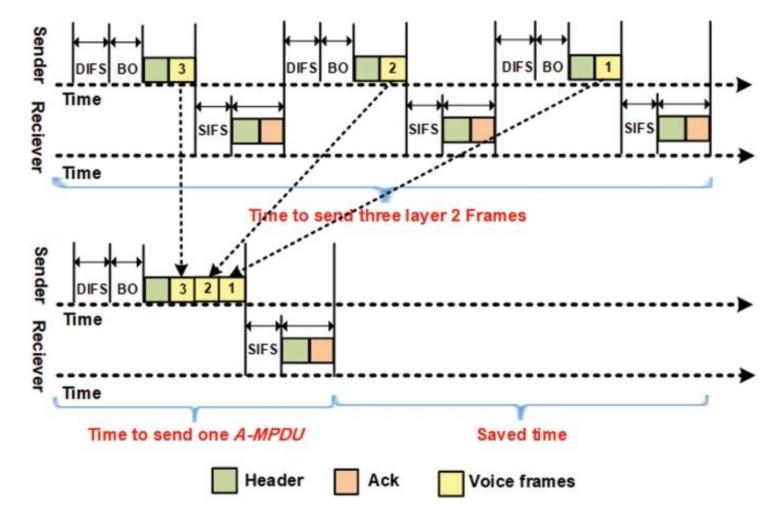




Frame Aggregation



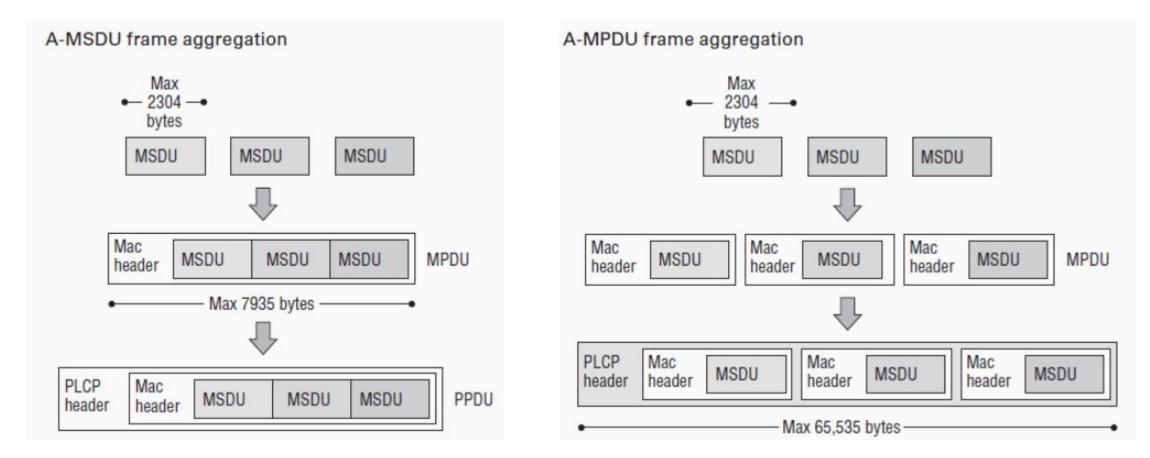
Frame aggregation increases throughput by sending multiple data frames in a single transmission. It reduces 802.11 protocol overhead, as multiple packets can be sent with a single PHY and MAC header, instead of each packet having its own headers. The number of ACKs and interframe spaces (and contention periods, if not in a TXOP) is also reduced.





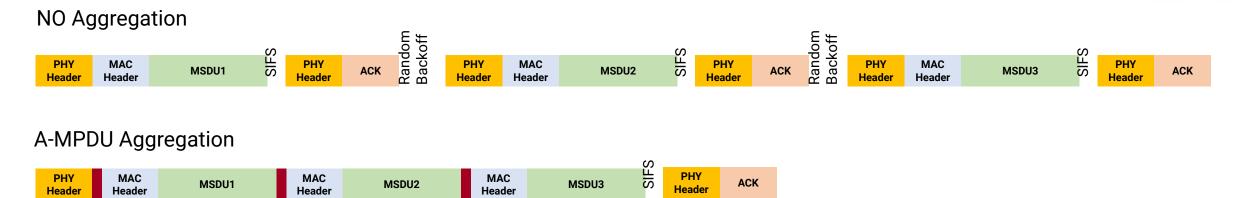
MAC Service Data Unit (MSDU) aggregation: the packets received by the MAC from the upper layer are MSDUs. Each packet gets an MSDU subframe header. Two or more subframes are bundled together and put in an 802.11 MAC frame (header + trailer). The resulting frame is an aggregate-MSDU (a-MSDU). The a-MSDUs are transmitted with a single PHY header by the radio.

MAC Protocol Data Unit (MPDU) aggregation: MPDUs are frames passed from the MAC to the PHY layer. Each MPDU has a MAC header and trailer. Multiple MPDU-s are bundled together to create an aggregate MPDU (a-MPDU), which is transmitted with a PHY header by the radio.



Types of Frame Aggregation





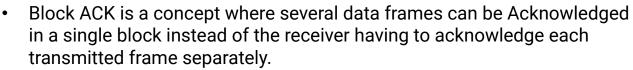
A-MSDU Aggregation

| ' | | S | | | | | | |
|---|---------------|---------------|-------|-------|-------|-----|---------------|-----|
| | PHY Header | MAC Header | MSDU1 | MSDU2 | MSDU3 | SIF | PHY Header | ACK |

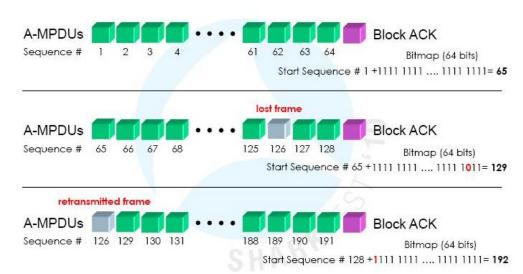
A-MSDUs inside A-MPDU Aggregation

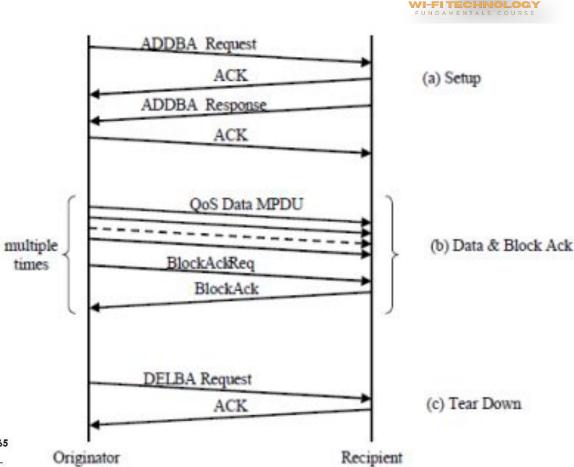
| | | | | | | | S | | |
|---------------|---------------|-------|-------|---------------|-------|-------|-----|---------------|-----|
| PHY Header | MAC Header | MSDU1 | MSDU2 | MAC Header | MSDU3 | MSDU4 | SIF | PHY Header | АСК |

Block Acknowledgements



- Block ACK is used during Aggregation.
- First Transmitter needs to check if the receiver supports block ACK feature.
- To find out TX can send Add Block ACK (ADDBA) request and to that the RX can respond with a ADDBA response.
- After data transmission is completed the TX can send a Block Ack Request Frame to received the Block ACK.
- The block ACK frame will have bitmap that indicates which frames were not received and this allows the Tx to retransmit those frames.
- At the end of the session the Tx can close by sending a Delete Block Ack (DELBA) message

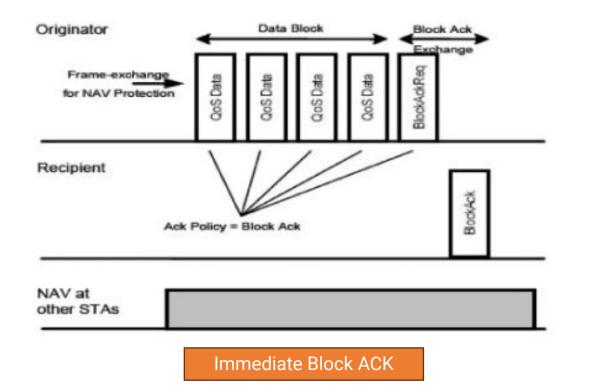


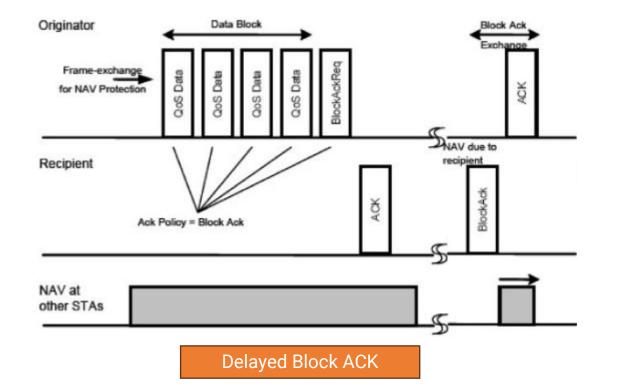


Immediate vs Delayed Block ACK



- There are two types of Block Ack mechanisms: immediate and delayed.
- Immediate Block Ack is suitable for high-bandwidth, low latency traffic
- Delayed Block Ack is suitable for applications that can tolerate moderate latency





Effects of Aggregation on Throughput

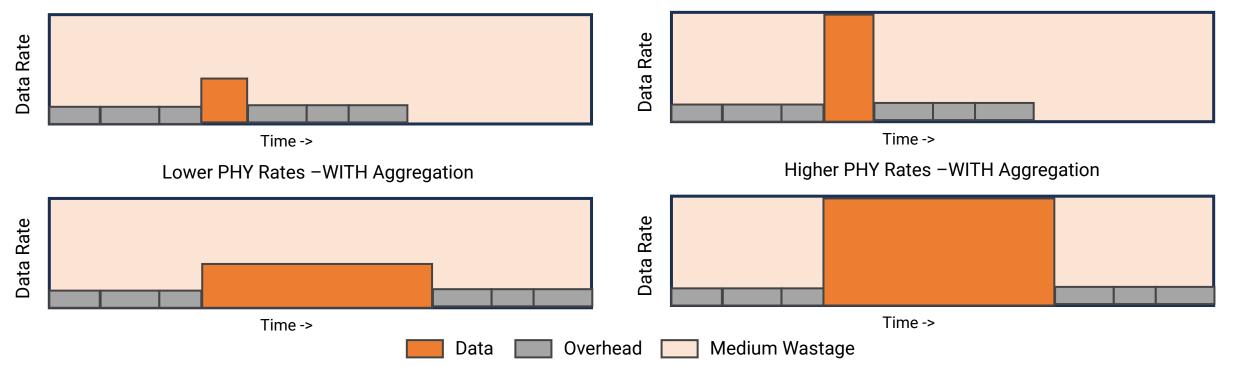


- At lower data rates the difference is small but at high data rates, there is a huge difference in throughput with and without aggregation.
- At smaller packet sizes also there is not much difference but at large packet sizes the difference is very big.

| HT Mixed Mode 1500 Bytes | | WITHOU | T Aggregation | WITH Aggregation | | |
|--------------------------|---------------------------------|------------|------------------|------------------|------------------|--|
| PHY Type | PHY Rate | Throughput | % Medium Wastage | Throughput | % Medium Wastage | |
| 11n 20MHz 1SS | 72.2Mbps | 32.45 Mbps | 55% | 67.76 Mbps | 6% | |
| 11n 20MHz 2SS | 144.4Mbps | 42 Mbps | 71% | 132.99 Mbps | 8% | |
| 11n 20MHz 3SS | 216.7Mbps | 45.16 Mbps | 79% | 193.71 Mbps | 11% | |
| 11n 20MHz 4SS | 288.9Mbps | 48.05 Mbps | 83% | 251.28 Mbps | 13% | |
| 11n 40MHz 1SS | 150Mbps | 43.21 Mbps | 71% | 137.94 Mbps | 8% | |
| 11n 40MHz 4SS | 1n 40MHz 4SS 600Mbps 52.24 Mbps | | 91% | 469.4 Mbps | 21% | |

Lower PHY Rates – WITHOUT Aggregation

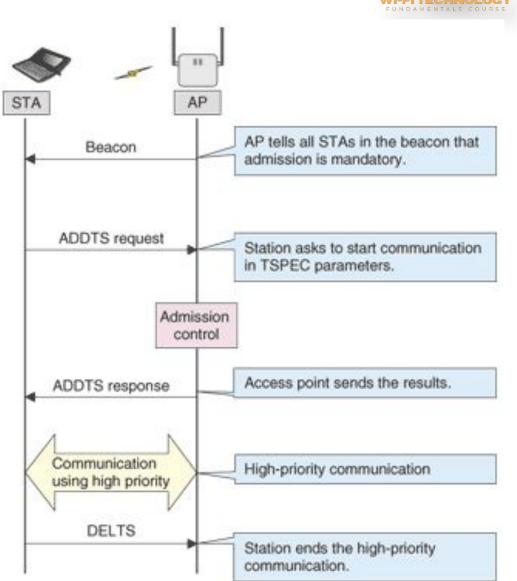
Higher PHY Rates –WITHOUT Aggregation



802.11e Admission Control

- The TSPEC (traffic specifications) negotiation procedure is defined in the IEEE802.11e standard
- Station asks the AP for its QoS requirements, such as mean data rate, packet length, and physical rate via an add traffic stream (ADDTS) request.
- The AP decides whether the request is acceptable or not and transmits its decision to the station.
- The station can start high-priority communication only when it is permitted to do so by the AP.
- The station also sends a delete traffic stream (DELTS) message when it has finished communicating.

TSPEC negotiation can prevent the wireless link from becoming congested and can keep the communication quality good. Stations then know that congestion has occurred in the wireless link before they start communicating and can wait to connect. This enables stations to use real-time applications such as VoIP and video comfortably.



Source: https://www.ntt-review.jp/archive/ntttechnical.php?contents=ntr200711sf7.html



QBSS Element

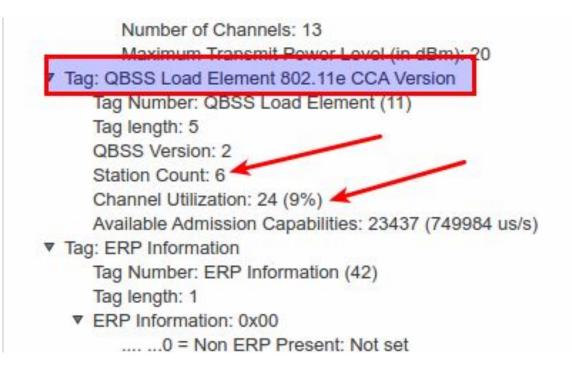


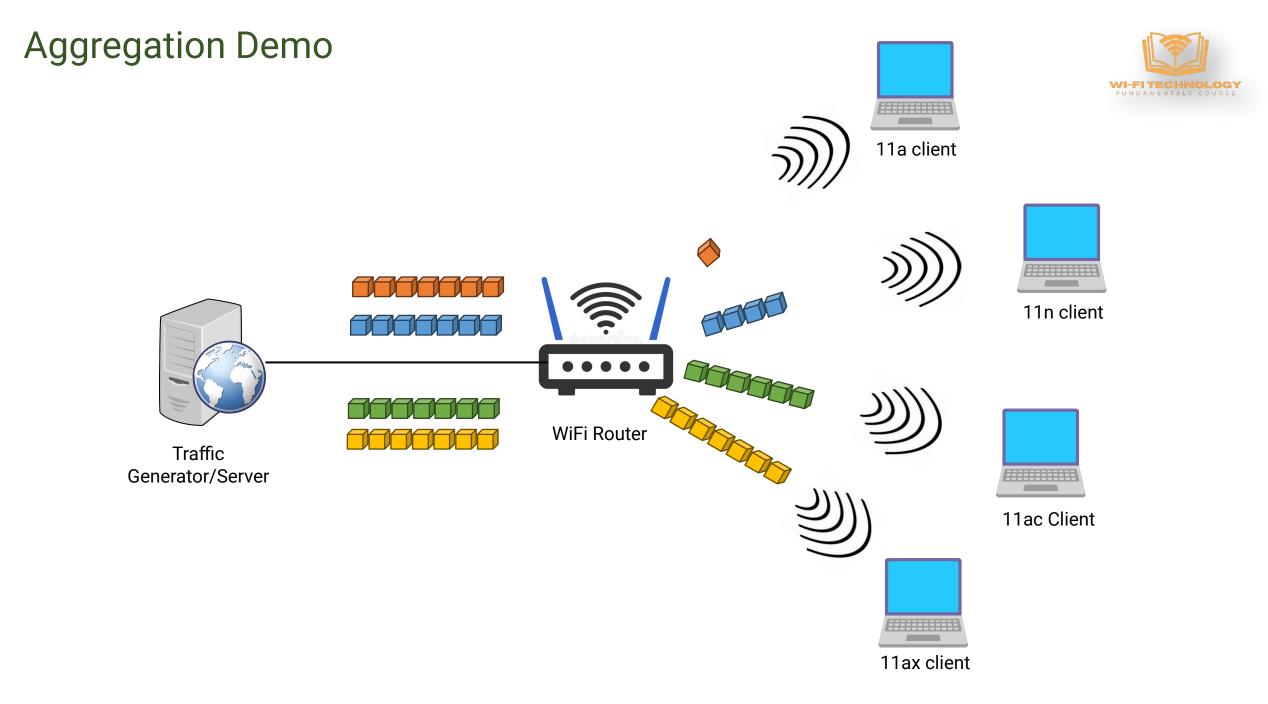
The QBSS (QOS enhanced basic service set) information element is an 802.11e construct that enables an access point to communicate its channel usage to wireless clients. It is intended to solve the problem of candidate access point selection, or intelligent roaming. It is located in the beacon frames of access points.

Station Count : Indicates the total number STAs currently associated with the QBSS (2 octets, unsigned integer value)

Channel Utilization: A percentage of time (normalized to 255) that the QAP sensed the medium was busy, Medium busy measured as physical or virtual carrier sense (CS) mechanism

Available Admission Capacity: Signals the remaining amount of medium time available via explicit admission control, Values from 0-31250 (2 octets long), Units of 32 μ s/s









Wi-Fi Airtime Calculator https://gjermundraaen.com/thewifiairtimecalculator/

Admission and Traffic Control Techniques in WLANs <u>https://www.ntt-review.jp/archive/ntttechnical.php?contents=ntr200711sf7.html</u>

802.11 Aggregation https://www.youtube.com/watch?v=3jqYwFQSqnE

Evaluation of IEEE 802.11 coexistence in WLAN deployments <u>https://typeset.io/papers/evaluation-of-ieee-802-11-coexistence-in-wlan-deployments-yvgb63w3f4</u>



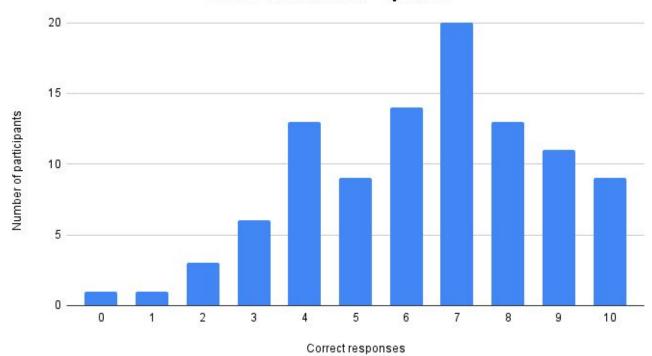




Quiz 3c Results



Number of participants - 100



Score distribution - quiz 3c

Winner Madhu R INDIA